A massively scalable distributed multigrid framework for nonlinear marine hydrodynamics

Stefan Lemvig Glimberg¹, Allan Peter Engsig-Karup¹,² and Luke N Olson³

Abstract
The focus of this article is on the parallel scalability of a distributed multigrid framework, known as the DTU Compute GPUlab Library, for execution on graphics processing unit (GPU)-accelerated supercomputers. We demonstrate near-ideal weak scalability for a high-order fully nonlinear potential flow (FNPF) time domain model on the Oak Ridge Titan supercomputer, which is equipped with a large number of many-core CPU-GPU nodes. The high-order finite difference scheme for the solver is implemented to expose data locality and scalability, and the linear Laplace solver is based on an iterative multilevel preconditioned defect correction method designed for high-throughput processing and massive parallelism. In this work, the FNPF discretization is based on a multi-block discretization that allows for large-scale simulations. In this setup, each grid block is based on a logically structured mesh with support for curvilinear representation of horizontal block boundaries to allow for an accurate representation of geometric features such as surface-piercing bottom-mounted structures—for example, mono-pile foundations as demonstrated. Unprecedented performance and scalability results are presented for a system of equations that is historically known as being too expensive to solve in practical applications. A novel feature of the potential flow model is demonstrated, being that a modest number of multigrid restrictions is sufficient for fast convergence, improving overall parallel scalability as the coarse grid problem diminishes. In the numerical benchmarks presented, we demonstrate using 8192 modern Nvidia GPUs enabling large-scale and high-resolution nonlinear marine hydrodynamics applications.

Keywords
High-performance computing, multi-GPU, domain decomposition, Laplace problem, multi-block solver, geometric multigrid, heterogeneous computing, free surface water waves

I. Introduction
The objective of this work is to detail and benchmark a newly developed distributed multigrid framework, which we refer to as the DTU Compute GPUlab Library (Glimberg, 2013). The framework offers scalable execution on supercomputers and compute cluster with heterogeneous architectures equipped with many-core coprocessors such as graphics processing units (GPUs). The GPUlab library builds upon thrust (Bell and Hoberock, 2011) and the message passing interface (MPI) (Gropp et al., 1999) and is designed for extensibility and portability, allowing for fast prototyping of advanced numerical solvers. In addition, using a layer of abstraction, the library delivers an accessible interface, allowing developers to realize high-performance through auto-tuning. The generic design principles used for implementing the library components are described in Glimberg et al. (2013).

To target marine hydrodynamic simulations, we have developed a GPUlab module based on fully nonlinear potential flow (FNPF) theory for free surface flows. This module is referred to as OceanWave3-D-GPU and is developed to achieve $O(n)$-scaling with $n$ degrees of freedom in

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the discretization. To achieve scalability, we utilize a massively parallel multilevel preconditioned defect correction scheme proposed in Engsig-Karup et al. (2011). In this article, we detail several extensions including the use of a multi-block procedure to enable multi-GPU computing and the development of boundary-fitted curvilinear domains to support real hydrodynamics application. Together, the OceanWave3-D-GPU solver is capable of targeting large marine areas due to the high parallel performance and low memory footprint in the simulation.

The OceanWave3-D-GPU simulator is an important step in expanding the concept of numerical wave tanks (NWTs) (Nimmala et al., 2013) toward more complex marine hydrodynamic applications using FNPF models. Fast hydrodynamics codes have shown potential in real-time naval hydrodynamics simulations and visualization (Engsig-Karup et al., 2011; Glimberg, 2013; Lindberg et al., 2013); efficient uncertainty quantification (Bigoni et al., 2016); tsunami propagation (Grilli et al., 2002); and the development, design, and analysis of marine structures such as naval vessels, wave-energy conversion devices (Verbrugghe et al., 2016), and coastal infrastructures. The multitude of applications motivates the use of both nonlinear wave–wave and wave–structure interactions.

1.1. Applications for large-scale water wave simulations

Accurate description of the nonlinear evolution and transformation of ocean waves over possibly changing sea floors is important in the context of ocean modeling and engineering. Recent advances in scalable numerical algorithms allow us to consider fully nonlinear and dispersive wave equations based on full potential theory in the context of large-scale simulations, for which computational fluid dynamics (CFD)-based models are too computationally expensive (Hino et al., 2014). The immediate advantage is a more complete description in the model and a wider applicability for hydrodynamics calculations, though still subject to the assumptions of non-breaking and irrotational waves. However, model extensions to account for breaking waves are possible (Kazolea and Ricchiuto, 2016). Simulations based on the fully nonlinear and dispersive equations are attractive for engineering applications where water waves transform nonlinearly over varying bathymetry and when accurate kinematics are required—for example, the estimation of structural loads. With a large-scale model, wave–wave interactions can be modeled even in the presence of multiple structures covering large areas, such as offshore wind farms.

A comprehensive introduction to hydrodynamics is given in Svendsen and Jonsson (1976), while details on numerical modeling are found in Lin (2008) or in the shorter survey by Dias and Bridges (2006). A variety of numerical techniques for solving such hydrodynamic model problems have been presented throughout the literature—for example, see Cai et al. (2005), Engsig-Karup (2006), Li and Fleming (1997), and Yeung (1982). Finally, a concise derivation of the fully nonlinear water wave model based on potential flow theory is revisited in Engsig-Karup et al. (2013).

1.2. Coprocessors

High-performance computing has been transformed over the last decade with the emergence of programmable many-core coprocessors, such as GPUs. This is driving the need for developing software and algorithms that exploit the parallelism offered by their compute units. Models based on full potential theory have previously used a variety of discretization methods. Due to the memory layout and access efficiency on a GPU, the finite difference method has been used to efficiently solve the governing potential flow equations. Memory bandwidth is the performance bottleneck for most scientific computing applications, and efficient memory access is therefore key for overall efficiency; this is also supported by recent trends in hardware development (Asanovic et al., 2006; Holesinski et al., 2012; Micikevicius, 2009).

In recent work, we have addressed aspects of robust, accurate, and fast three-dimensional (3-D) simulation of water waves over uneven sea floors (Engsig-Karup et al., 2008, 2011; Glimberg et al., 2011) based on full potential theory. The objective of this work is to enable new hydrodynamic applications by extending the scalability on distributed systems for use with a large number of time steps and billions of spatial degrees of freedom. To this end, we present and benchmark an extension of the parallel model (Engsig-Karup et al., 2011) that supports large GPU-accelerated systems using a hybrid MPI-CUDA implementation.

This work is related to the methods originally proposed in Li and Fleming (1997, 2001), which is based on multi-grid using low-order discretization. This was extended to a high-order finite difference technique in two-dimensional (2-D) (Bingham and Zhang, 2007) and to 3-D (Engsig-Karup et al., 2008) approaches, and the approach enables efficient, scalable, and low-storage solution of the Laplace problem using an efficient multigrid strategy proposed and analyzed in Engsig-Karup et al. (2011) and Engsig-Karup (2014).

A novel feature of our approach is that communication requirements are minimized due to a low number of multigrid restrictions, leading to low overhead on massively parallel machines and high scalability. Our benchmark examples verify that the combination of a strong vertical smoother and a favorable coarsening strategy based on aspect ratios provides good convergence at even few multigrid levels. Reducing the number of multigrid levels leads to a lower communication overhead.

1.3. On GPU-accelerated multigrid solvers for scientific applications

The use of coprocessors, such as GPUs, for elliptic problems was first reported in Fan et al. (2004). In the setting of
finite elements, weak scalability is explored in Goddeke et al. (2007) and later on clusters in Goddeke (2011). Additional research on both geometric and algebraic multigrid methods have been demonstrated on a number of GPU-accelerated applications in various engineering disciplines (Bell et al., 2012; Esler et al., 2014; Naumov et al., 2015; Strzodka et al., 2013). The increasing ratio between compute performance and memory bandwidth favors higher order discretizations that increase the flop counts to memory transfers and make it possible to tune trade-offs between accuracy and numerical efficiency.

1.4. Paper contributions

We present a novel and scalable algorithmic implementation for the FNPF model enabling large-scale marine hydrodynamics simulation in realistic engineering applications. The implementation extends the proof-of-concept for a single-GPU system demonstrated in Engsig-Karup et al. (2011). This is facilitated by a multigrid framework described in Glimberg (2013) that is based on an MPI-CUDA implementation. The framework handles curvilinear and logically structured meshes to introduce geometric flexibility—for example, the simulation of mono-pile foundations in marine areas. Through a careful semi-coarsening strategy, we demonstrate that the rapid convergence of the iterative solver requires only a few multigrid levels, avoiding an expensive coarse grid problem that typically limits scalability. We provide a benchmark and demonstrate weak scalability on several compute clusters, including the Oak Ridge Titan supercomputer, utilizing up to 8192 GPUs and enabling high-resolution and large-scale marine hydrodynamic simulation.

2. The mathematical model

In recent works, we have focused on the robustness of the numerical scheme (Engsig-Karup, 2014; Engsig-Karup et al., 2008; Li and Fleming, 1997) along with efficiency on emerging many-core architectures (Engsig-Karup et al., 2011; Glimberg et al., 2011). A mathematical model for potential free surface flow in three spatial dimensions can be expressed in terms of the unsteady kinematic and dynamic boundary equations in Zakharov form

\[ \partial \eta = -\nabla \eta \cdot \nabla \phi + \tilde{w}(1 + \nabla \eta \cdot \nabla \eta) \]  
\[ \partial \phi = -g\eta - \frac{1}{2} \nabla \phi \cdot \nabla \phi - \frac{\tilde{w}^2}{2} (1 + \nabla \eta \cdot \nabla \eta) \] (2)

where \( \nabla \equiv [\partial_x, \partial_y, \partial_z]^T \) is the horizontal gradient operator, \( g \) is the gravitational acceleration, and \( \eta(x,t) \) is the free surface elevation measured from the still water level at \( z = 0 \). In addition, \( \hat{\phi}(x,t) \) is the scalar velocity potential, and \( \tilde{w}(x,t) \) is the vertical velocity, both are evaluated at the free surface \( z = \eta \). Assuming that the flow is inviscid, irrotational, and non-breaking, the continuity equation can be expressed in terms of a scalar velocity potential function in the form of

\[ [\nabla^2 + \partial_z] \phi = 0 \] (3)

which requires a closure for the free surface equations. By imposing the free surface boundary conditions (2) together with a kinematic bottom boundary condition

\[ \partial_z \phi + \nabla h \cdot \nabla \phi = 0, \quad z = -h \] (4)

and assuming impermeable vertical walls at domain boundaries, the resulting Laplace problem is well-posed. From the solution, the velocity field in the fluid volume is \( (u, w) = (\nabla \phi, \partial_z \phi) \).

2.1. Vertical coordinate transformation

To improve efficiency, a fixed time-invariant computational domain is introduced using a classical \( \sigma \)-transformation

\[ \sigma(x,z,t) = \frac{z + h(x,t)}{h(x) + \eta(x,t)} \] (5)

which maps \( z \rightarrow \sigma \) and where \( -h \leq z \leq \eta \) and \( 0 \leq \sigma \leq 1 \). The resulting transformed Laplace problem with \( \phi(x,z) = \Phi(x,\sigma) \) is written as

\[ \sigma = 1:\]

\[ \Phi = \hat{\phi} \] (6)

\[ 0 \leq \sigma < 1:\]

\[ \nabla^2 \Phi + (\nabla^2 \sigma)(\partial_z \Phi) + 2 \nabla \sigma \cdot \nabla (\partial_z \Phi) + (\nabla \sigma \cdot \nabla + (\partial_\sigma)^2) \partial_{zz} \Phi = 0 \] (7)

\[ \sigma = 0:\]

\[ (\partial_z \sigma + \nabla h \cdot \nabla \sigma)(\partial_z \Phi) + \nabla h \cdot \nabla \Phi = 0. \] (8)

These equations, together with the kinematic and dynamic free surface boundary conditions, describe the FNPF model in Cartesian coordinates, applicable to applications where the physical domain can be represented within an NWT with regular boundaries.

2.2. Horizontal mapping to generalized coordinates

Irregular horizontal domains can be introduced by rewriting the Cartesian coordinates and by expressing the physical derivatives in terms of the computational domain. This is accomplished by introducing the affine coordinate mapping \( \chi(x,y) = \{(x,y) \rightarrow (\xi, \gamma)\} \). Then, applying the chain rule, the first-order derivatives with respect to the original physical coordinates \((x,y)\) of a function \(u\) are described within the regular computational domain \((\xi, \gamma)\) with the following relations (subscripts are used to denote derivatives)

\[ u_x = \xi_x u_{\xi} + \gamma_x u, \] (9)

\[ u_y = \xi_y u_{\xi} + \gamma_y u, \] (10)
under which the following relations hold
\[
\xi_x = \frac{1}{J} y_y, \quad \xi_y = -\frac{1}{J} x_x \tag{11}
\]
\[
\gamma_x = -\frac{1}{J} y_z, \quad \gamma_y = \frac{1}{J} x_z \tag{12}
\]
where \( J \) is the Jacobian of the mapping
\[
J = \det(J) = x_y y_z - x_y z_y \tag{13}
\]
Given the above equations, the first-order derivatives of \( u \) can be described exclusively within the computational reference domain as
\[
u_x = \frac{y_x u_x - y_x u_y}{J} \quad \tag{14}
\]
\[
u_y = \frac{x_x u_x - x_x u_y}{J} \quad \tag{15}
\]
In a similar way, the second-order and mixed derivatives can be derived. Under the assumption of time-invariant domains, the Jacobian and its derivatives are constant and may be precomputed to improve computational efficiency.

Having assumed impermeable wall boundaries, the net-flux through these boundaries is zero and is stated in terms of a homogeneous Neumann boundary condition
\[
\mathbf{n} \cdot \nabla u = 0, \quad (x, y) \in \partial \Omega \tag{16}
\]
where \( \mathbf{n} = (n_x, n_y)^T \) is the outward pointing horizontal normal vector defined at the vertical boundary \( \partial \Omega \) surrounding the domain \( \Omega \). The normal vector in the physical and generalized coordinates is connected through the following relation
\[
\begin{pmatrix} n_x \\ n_y \end{pmatrix} = \begin{pmatrix} \xi_x & -\xi_y \\ \sqrt{\xi_x^2 + \xi_y^2} & \sqrt{\xi_x^2 + \xi_y^2} \\ \gamma_x & -\gamma_y \\ \sqrt{\gamma_x^2 + \gamma_y^2} & \sqrt{\gamma_x^2 + \gamma_y^2} \end{pmatrix} \begin{pmatrix} n_x \\ n_y \end{pmatrix} \tag{17}
\]
Note that the normals in the computational domain are trivial, as they are always perpendicular to the rectangular domain boundaries.

Supporting curvilinear coordinates enables complex modeling of offshore structures, harbors, shorelines, and so on, which have significant engineering value over traditional regular wave tank domains (Fang et al., 2012; Li and Zhan, 2001; Shi and Sun, 1995; Shi et al., 2001; Zhang et al., 2005). Flexible representation of the discrete domain can also be utilized to spatially adapt the grid to the wavelengths, resulting in higher resolution where shorter waves are present and thereby minimizing over- or under-resolved waves.

3. The numerical model
All spatial derivatives in equations (1), (2), and (6) to (8) are discretized using standard centered finite differences of flexible-order accuracy, meaning that the size of the stencil operators can be pre-adjusted to allow different orders of accuracy using the method of undetermined coefficients. The stencil implementation is matrix-free to avoid the use of sparse matrix formats, which would require additional index lookups from memory. The stencil operations also allow efficient access to GPU memory as data can be cached (Micikevicius, 2009). The temporal surface conditions in equations (1) and (2) are solved with a fourth-order explicit Runge–Kutta method. The time domain \([0, T_{end}]\) is partitioned uniformly into discrete steps \( \Delta t = \Delta t_0 \), \( \Delta t = 1, 2, \ldots \), such that \( \Delta t \) satisfies the courant–friedrichs–lewyl (CFL) stability condition. At each stage of the Runge–Kutta method, the vertical velocity at the surface, \( \tilde{w} \), is approximated as the vertical derivative of the velocity potential \( \Phi \), where \( \Phi \) is described by the linear system of equations due to the discretization of equation (5)
\[
\mathcal{A} \Phi = \mathbf{b}, \quad \mathcal{A} \in \mathbb{R}^{N \times N}, \mathbf{b} \in \mathbb{R}^{N \times 1} \tag{18}
\]
where \( N \) is the total degree of freedom. The sparse, nonsymmetric system (18) is a key computational bottleneck in the simulation and is the focus of our attention. To extend the original work by Li and Fleming (1997), a multigrid preconditioned defect correction method can be employed, which has been expanded for improved accuracy, efficiency, and robustness (Bingham and Zhang, 2007; Engsig-Karup et al., 2008, 2011). The iterative defect correction method is attractive because it has a short and highly parallel outer loop, requiring only one collective all-to-one communication step when computing the residual norm for evaluating the convergence criteria. The preconditioned defect correction method generates a sequence of iterates \( \mathbf{f} \), starting from an initial guess \( \Phi^0 \)
\[
\Phi^n = \Phi^{n-1} + \mathcal{M}^{-1} \mathbf{r}^{n-1}, \quad j = 1, 2, \ldots \tag{19}
\]
where \( \mathbf{r}^{n-1} = \mathbf{b} - \mathcal{A} \Phi^{n-1} \) denotes the residual at iteration \( j - 1 \) and \( \mathcal{M} \) is a left preconditioner based on nested grids to form a geometric multigrid preconditioner (Trottengen et al., 2000). The iterations continue until the convergence criteria are satisfied
\[
\| \mathbf{r}^j \| \leq rtol \| \mathbf{b} \| + atol \tag{20}
\]
where \( rtol \) and \( atol \) are the relative and absolute tolerances, respectively. Previous work has demonstrated that \( \mathcal{M}^{-1} \) is effective when based on linearized second-order (low-order) finite differences to allow efficient computation of the preconditioned system, while still obtaining fast algorithmic convergence to the original system of high-order accuracy. In addition, such low-order approximation significantly reduces the computational work of computing the residuals and relaxations during the preconditioning phase.

4. Spatial domain and data distribution
It is often advantageous to decompose a boundary value problem into smaller subdomains. For one, the
computational work can be distributed and solved in parallel to achieve better overall performance. In addition, memory distribution lowers the memory requirements per compute node and allows for larger global problems. However, communication between compute nodes requires frequent message passing, leading to degraded performance, especially for smaller problem sizes.

For the free surface water wave model, a data decomposition technique with small domain overlaps is proposed. The computational domain is first decomposed into rectangular nonoverlapping subdomains, and then the so-called ghost layers are added to account for grid points that are distributed across adjacent subdomains (Acklam and Langtangen, 1998). The size of the ghost layers depends on the size of the finite difference stencil. An example of the decomposition of a simple $17 \times 5$ domain into two subdomains is illustrated in Figure 1. Ghost points are updated block-wise, indicated by the arrows, when information from adjacent domains are queried. Thus, subdomains concurrently solve the global boundary value problem by communicating with neighboring subdomains. This approach differs from domain decomposition methods where overlapping local boundary value problems are solved individually (Smith et al., 1996). Our proposed approach requires all ghost layers to be updated prior to any global operation, but most importantly it maintains the attractive convergence rate of the multigrid preconditioned defect correction method.

The vertical resolution for any free surface model is significantly lower than the horizontal resolution. Ten vertical grid points, for example, are often sufficient for hydrodynamics, while the number of horizontal grid points is dictated by the size of the physical domain and the wavelengths. The global domain is, therefore, decomposed using a 2-D horizontal topology. For practical applications, the necessary spatial resolution is typically less than 10 nodes in the vertical direction leading to sufficient accuracy in the dispersion relation for wave propagation. Thus, for applications, a large number of nodes can be used to represent the horizontal dimensions covering large marine areas, for example, to simulate long wave propagation in coastal and harbor areas with varying bathymetry. Such large domains will have a favorable ratio between internal grid points and ghost points compared to smaller domains, causing less communication overhead relative to the computation volume.

### 4.1. Multi-GPU data distribution

A multigrid preconditioner, $M$, can be formed as a series of (sparse) matrix–vector operations. Any such matrix–vector operation, $S$, acting on a subdomain $\Omega_i$, may be formulated as

$$S_i \left( x_i, g_i \right) = b_i, \quad S_i \in \mathbb{R}^{n \times n}, b_i \in \mathbb{R}^{n \times 1}$$

where $n$ is number of degrees of freedom in the subdomain, requiring $n \leq N$. The vectors $x_i$ and $g_i$ contain the internal points and the ghost points of subdomain $i$, respectively. Computing the matrix–vector product using $S_i$ is a local operation, which can be performed by the process (GPU) associated with subdomain $\Omega_i$. To facilitate this, all ghost layers are updated prior to applying $S_i$ by transferring boundary layer data between GPUs using MPI. The communication overhead due to the exchange of ghost values is related to the size of the ghost layers $g_i$ and is investigated in the following sections. When updating the internal boundaries between neighboring GPUs, memory is first copied from the GPU to the CPU before transferred via MPI. To reduce transfer overhead, this is performed first for the east/west boundaries, to allow asynchronous MPI communication, while the north/south boundaries are copied to the CPU. Figure 2 depicts some typical scenarios. The single workstation offers low memory transfer overhead as no network communication is required. Yet the number of GPUs on a single workstation is limited due to the number of PCIe sockets (often one or two). In contrast, on a cluster or supercomputer, a distributed memory model is provided, allowing a flexible number of nodes for simultaneous computations, at the expense of lower network bandwidth.

5. Multi-GPU performance study

The numerical method has been verified and compared to both analytic and experimental data in previous works (Engsig-Karup et al., 2011, 2013). One advantage of the multi-GPU decomposition method presented here is that it preserves the attractive algorithmic efficiency of a single-block multigrid method, which can be confirmed by verifying that the residual norms for both the single- and multi-GPU implementations agree within machine precision. In the following, $\Omega_i$ represents the $i$th subdomain, with $i = 1, \ldots, P$, where $P$ is the number of subdomains. The number of subdomains is not restricted to match the number of GPUs; however, no performance gains are expected to compare a one GPU per subdomain. In the following sections, the number of subdomains and GPUs are, therefore, always equal.

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**Figure 1.** A decomposition of a grid of global size $17 \times 5$ into two subdomains with two layers of ghost points. ● and ■ represent internal grid points and ghost points, respectively. ○ indicates internal points that may be shared between grids to ensure an uneven dimension required by the geometric multigrid solver.
The following discussion focuses on the numerical efficiency and convergence of the multi-GPU implementation for solving the Laplace problem (18) for varying problem sizes and increasing number of GPUs. A detailed performance analysis of an optimized single-GPU version can be found in Engsig-Karup et al. (2011). A desktop computer and three compute clusters have been used to study the multi-GPU performance. Machine details are summarized in Table 1.

5.1. Data distribution performance

The computational bottleneck is solving the Laplace problem. Introducing multiple subdomains and thereby multiple GPUs influences the performance of each individual subroutine of the multigrid preconditioned defect correction algorithm. As a measure of performance, the average timings based on 100 defect correction iterations are recorded for an increasing number of unknowns. One multigrid V-cycle with Red-Black Z-line relaxation is used for preconditioning together with sixth-order accurate finite differences for the spatial discretization. The number of coarsening levels in a V-cycle is denoted by $K$. The timings in Table 2 use $K = K_{\text{max}}$ for each run, meaning that each V-cycle continues until a $5 \times 5 \times 3$ grid size per subdomain. A semi-coarsening strategy that best preserves the grid isotropy is chosen, implying that the grid is restricted only in the horizontal dimensions until the grid spacing is of similar size as the vertical grid spacing. Hereafter, all dimensions are restricted. At each multigrid level, two pre- and post-relaxations are used along with four relaxations at the coarsest level. The relative and absolute timings for each subroutine and for an increasing number of subdomains, $P$, are reported in Table 2. Residual (high) in the first column refers to the sixth-order residual evaluation in the defect correction loop, while residual (low) refers to the second-order linear residual evaluation in the preconditioning phase.

Note from the numbers in Table 2 that increasing the number of subdomains, and thereby the number of GPUs improves the overall computational time only for problems larger than $513 \times 513 \times 9$. We believe this is acceptable, as it is often difficult for multiple GPUs to efficiently solve problem sizes that fit within the memory of a single GPU. Since each GPU is massively parallel, using multiple GPUs for a small problem introduces costly overhead, which the current implementation has not been able to overcome. Strong scaling (in terms of number of GPUs) for problems of moderate sizes should, therefore, not be expected to be good in general.

Based on the relative timings in Figure 3 we see that relaxation becomes increasingly dominant as the number of subdomains (thus GPUs) increases. Communication between GPUs occurs at every relaxation step, which happens multiple times at every multigrid level. For the coarse grid levels, the surface to volume ratio is larger and hence communication overhead becomes a dominating performance bottleneck. It is, therefore, particularly desirable to decrease the number of multigrid levels to also reduce the number of relaxation and consequently lower communication requirements. However, multigrid achieves its unique algorithmic scalability and grid-independent convergence properties from the fact that it reduces error frequencies by relaxing on all grid levels. This is in general important for elliptic problems where all grid points are coupled.

In Section 5.3., a test example illustrates the effect of grid levels and horizontal to vertical aspect ratio on algorithmic and numerical efficiency.

5.2. Numerical performance of multigrid restrictions

As a consequence of the previous observations, an examination of how the number of multigrid restrictions affects the numerical performance for both the single and the multi-block solver is performed. The absolute time per outer defect correction iteration is used here as a measure for performance. Thus, timings are independent of any physical properties of the free surface problem since algebraic convergence is not considered. Timings are measured and reported in Table 3 for a variation of levels $K$ and number of subdomains $P$. The remainder of the solver properties is the same as in the previous example.

Two different speedup measures are reported in Table 3: $\delta_K$ is the speedup for $K$ levels in comparison with $K_{\text{max}}$, with a fixed number of subdomains $P$. Likewise, $\delta_P$ is the...
speedup for $P$ subdomains in comparison with one domain, with a fixed number of restrictions $K$. We see that $d_K = 1/21$, $d_K = 3/21$, $d_K = 5/21$, and $d_{\text{max}}/21$ as expected. $d_P$ is a measure of strong scaling with respect to the number of GPUs (subdomains). As mentioned earlier, multiple GPUs are feasible only for problems of reasonable sizes. Based on the reported numbers, we conclude that there is good potential for improving the overall numerical performance, given that few restrictions (low $K$) are sufficient for rapid convergence. In the following, a test is set up to demonstrate how this can be achieved.

### 5.3. Algorithmic performance of multigrid restrictions

To unify the findings from the previous examples, we now introduce and solve a specific nonlinear free surface problem. The purpose of this numerical experiment is to verify whether imposing fewer restrictions and thus fewer relaxations to minimize communication, based on the findings in Table 2, can in fact lead to performance improvements as reported in Table 3. The final link that needs to be demonstrated is whether the algebraic convergence can be maintained using fewer multigrid restrictions.

A 2-D Gaussian distribution is used as initial condition to the free surface elevation within an NWT with no flux boundaries:

$$h(x, y, t) = k e^{-x^2 + y^2 / 2}$$

where $k = 0.15$ and $\rho = 0.05$. The wavelength is approximately $\lambda = 1$ m and the wave number $k = 2\pi$. The distance from the seabed to still water level is $h = 1$ and, therefore, $kh = 2\pi$, which is intermediate depths. The physical size of the domain is extended to fit the initial aspect

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**Table 1.** Hardware configurations for the desktop computer and three large compute clusters used in the tests.

<table>
<thead>
<tr>
<th>Machine</th>
<th>Desktop</th>
<th>Oscar</th>
<th>Stampede</th>
<th>Titan</th>
</tr>
</thead>
<tbody>
<tr>
<td>Nodes (used)</td>
<td>1</td>
<td>16</td>
<td>32</td>
<td>8,192</td>
</tr>
<tr>
<td>CPU</td>
<td>Intel Xeon E5620</td>
<td>Intel Xeon 5540</td>
<td>Intel Xeon E5</td>
<td>AMD Opteron 6274</td>
</tr>
<tr>
<td>CPU frequency</td>
<td>2.40 GHz</td>
<td>2.53 GHz</td>
<td>2.4 GHz</td>
<td>2.2 GHz</td>
</tr>
<tr>
<td>CPU memory</td>
<td>12 GB</td>
<td>24 GB</td>
<td>32 GB</td>
<td>32 GB</td>
</tr>
<tr>
<td>GPU</td>
<td>4 × GeForce GTX 590</td>
<td>2 × Tesla M2050</td>
<td>Tesla K20m</td>
<td>Tesla K20X</td>
</tr>
<tr>
<td>CUDA cores</td>
<td>512</td>
<td>448</td>
<td>2496</td>
<td>2688</td>
</tr>
<tr>
<td>GPU Memory</td>
<td>1.5 GB</td>
<td>3 GB</td>
<td>5 GB</td>
<td>6 GB</td>
</tr>
</tbody>
</table>

GPU: graphics processing unit.

CUDA cores and memory are per GPU.

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**Table 2.** Relative and absolute timings for one defect correction iteration, divided into each subroutine, for an increasing number of unknowns and subdomains.

<table>
<thead>
<tr>
<th>Subroutine</th>
<th>$P$</th>
<th>$129 \times 129 \times 9$</th>
<th>$257 \times 257 \times 9$</th>
<th>$513 \times 513 \times 9$</th>
<th>$1025 \times 1025 \times 9$</th>
</tr>
</thead>
<tbody>
<tr>
<td>Residual (high)</td>
<td>1</td>
<td>15.2 0.0010</td>
<td>25.9 0.0036</td>
<td>34.2 0.0133</td>
<td>39.1 0.0529</td>
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*“Other” subroutines include an Axpy and a two-norm evaluation. Time in seconds on Oscar.*

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ratios that we wish to examine. The following number of grid points per wavelength is used $N_w = 4, 8, 16, 32$, which results in aspect ratios of $\Delta r/\Delta x = 4, 2, 1, 0.5$ at the finest grid level. Figure 4 illustrates how the nonlinear wave travels within the first seconds with different wave resolutions. The initial wave rapidly propagates throughout the domain and creates multiple waves of various amplitudes.

At each time stage, the Laplace problem (18) is solved to a relative and absolute tolerance of $10^{-4}$ and $10^{-5}$, respectively, then the number of outer defect corrections is counted. The simulation is continued until the average number of iterations for each defect correction does not change within the first three digits. The results are collected in Table 4 where the aspect ratios and number of restrictions are also reported. The results are encouraging as they indeed confirm that the number of useful restrictions is related to the discretization of the physical domain. Load balancing between the GPU threads and multiple GPUs will, therefore, not be a critical issue, as the discretization of most practical applications allows favorable aspect ratios requiring only few restrictions.

Since ghost layers are always updated prior to any operation, convergence is independent of the number of subdomains. Thus, the number of average iterations reported in Table 4 applies, regardless of the number of subdomains (GPUs) used for the same problem.

In addition to the above example, the spectral radius of the stationary iteration matrix for the two-level multigrid with Zebra-Line Gauss-Seidel smoothing is evaluated and plotted in Figure 5. A small spectral radius enables fast convergence, which is seen to be governed by the dimensionless variables for aspect ratio (vertical to horizontal) and wavelength ($L$) relative to ratio between seabed slope ($hx$) and depth ($h$) (Engsig-Karup, 2014; Engsig-Karup et al., 2008). As the semi-coarsening strategy seeks to keep $\Delta r/\Delta x$ close to 1 where the spectral radius is low, the figure also confirms that fast convergence can be achieved for wave propagation applications as intended.

![Figure 3](image)

**Figure 3.** Performance breakdown of each subroutine as the number of GPUs increases. Timings are for one preconditioned defect correction iteration for a problem size of $1025 \times 1025 \times 9$, on Oscar, compare, Table 2. GPU: graphics processing unit.

![Figure 4](image)

**Figure 4.** Nonlinear waves traveling in a closed basin with no flux boundaries, illustrated at three distinct time steps. Total horizontal grid dimensions are $129 \times 129$. (a) $T = 0$ s. (b) $T = 2$ s. (c) $T = 4$ s.
6. Large-scale performance analysis

Given the findings in previous sections, the solver performance is now evaluated on three larger GPU-accelerated clusters. Time per defect correction iteration is again used as the measure of performance since it is independent on the physical properties of the wave model. The waves are resolved with grid points such that three levels of the multigrid preconditioner are sufficient, $K = 3$, following the results from the previous section.

First performance test is performed in the Oscar cluster at Brown University. Oscar is equipped with Intel Xeon E5630 processors with 2.53 GHz and 32 GB host memory per compute node. Each node has two Nvidia Tesla M2050 GPUs. Absolute performance timings are summarized in Figure 6 as a function of increasing problem size. The single block timings ($\Omega_1$) evolve as expected, overhead of launching kernels are evident only for the smallest problem sizes, while for larger problems the time per iteration scales linearly. For the multi-GPU timings, a notable effect is observed when communication is dominating and when it is not. After approximately one million degrees of freedom, communication overhead becomes less significant and the use of multiple compute units starts to become beneficial. For the larger problem sizes, close to optimal linear performance is observed.

The second performance scale test is performed on the Stampede cluster at the University of Texas. Stampede is a Dell Linux cluster based on compute nodes equipped with two Intel Xeon E5 (Sandy Bridge) processors, 32 GB of host memory, and one Nvidia Tesla K20 m GPU. All nodes are connected with Mellanox FDR InfiniBand controllers. The Stampede cluster allows up to 32 GPU nodes to be

<table>
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<tr>
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</table>

GPU: graphics processing unit.

*a* The Avg. Solve columns indicate the total solve times on Oscar using one GPU, compare, the timings in Table 3.
occupied at once. Absolute performance timings are illustrated in Figure 6. Notice that a setup of 16 GPUs allows the solution of problems with more than one billion degrees of freedom in practical times. With an approximate time per iteration of $t_i = 0.6 \text{ s}$ at $N = 10^8$, it would be possible to compute a full time step in 3–6 s, assuming convergence in 5–10 iterations. With a time step size of $\Delta t = 0.05 \text{ s}$, a 1-min simulation can be computed in 1–2 h. This is well within a time frame considered practical for engineering purposes. With 32 GPUs available, compute time would reduce to almost half.

A final extremely large-scale performance evaluation is performed on the Titan cluster at the Oak Ridge National Laboratory. Titan nodes are equipped with AMD Opteron 6274 CPUs and Nvidia Tesla K20X GPUs. They have 32 GB host memory and are connected with Gemini interconnect. A scaling test up to 8192 GPUs is performed and reported in Figure 6(c). Timings are well aligned with the previous results, though the curves are less smooth. The implementation of the scalability test alternately doubles the number of grid points in the x- and y-direction. When increasing the problem in the x-direction, only coalesced memory at the ghost layers are introduced, whereas the y-direction adds uncoalesced ghost layers. Titan seems to be more sensible to such alternately coalesced/uncoalesced increases.

Weak and strong scaling results for all three compute clusters are depicted in Figure 7(a) and (b), respectively. Figure 7(a) shows the efficiency for each of the three compute clusters as the ratio between the number of GPUs and problem size is constant. There is a performance penalty when introducing multiple GPUs, indicated by the drop from one to multiple GPUs. Hereafter, weak scaling remains almost constant and there is no significant penalty when using additional GPUs. Weak scaling is stable up to the 8192 GPUs on Titan. The strong scaling figure shows the efficiency in terms of inverse cost for a given time consumption at a constant problem size of $N = 3.8 \times 10^7$. Ideally, increasing the number of GPUs would lead to a corresponding reduction in compute time, leading to constant curves in Figure 7(b). However, this is not the case for the given problem size, where parallel overhead becomes critical.

7. Vertical cylinder wave run-up

Predicting water scattering or wave loads on a vertical cylinder is relevant for offshore pile installations such as wind turbines. In this demonstration, the wave run-up around a single vertical cylinder in open water is considered using a boundary-fitted domain in curvilinear coordinates. Analytic formulations (MacCamy and Fuchs, 1954) and experimental data are available for linear and nonlinear incident waves over even seabed. For nonlinear water waves, experimental results are presented by Kriebel, 1990, 1992, where nonlinear diffraction effects are studied and determined to be of significant importance compared to linear theory. To demonstrate the application of a boundary-fitted domain, a nonlinear experiment from Kriebel (1992) is reproduced. Long and regular incident
waves are generated in the inlet zone with a wavelength of \( L = 2.73 \text{ m} \) and wave height \( H = 0.053 \text{ m} \). A wave damping outlet zone is used to avoid wave reflections (Glimberg, 2013). The still water depth \( h \) is constant and rather shallow, resulting in dimensionless \( kh = 1.036 \) and \( kH = 0.122 \). Diffraction occurs around the centered vertical cylinder of radius \( a = 0.1625 \), thus \( ka = 0.374 \). The nonlinear case uses a sixth-order finite difference approximation and a resolution of \( 1025 \times 129 \times 9 \). The vertical cylinder setup is presented in Figure 8.

Measuring the maximum fully developed wave run-up around the cylinder provides the \( 360^\circ \) wave height profile in Figure 9. \( 0^\circ \) is the opposite side of the pile relative to the incoming waves. The numerical results are well in agreement with the experiments, and they are consistent with previously presented numerical results using similar numerical approaches, for example, Ducroz et al. (2010). Although the vertical cylinder case is primarily an illustrative example, the introduction of an efficient solver for large-scale simulation, supporting generalized curvilinear coordinates, significantly increases the range of real engineering applications.

8. Conclusion and future work

For the Laplace problem of a fully nonlinear free surface model, we have presented a performance analysis for distributed parallel computations on cluster platforms. The spatial decomposition technique uses stencil-sized ghost layers, preserving the algorithmic efficiency and numerical properties of the established existing solvers for the problem, cf. recent studies of Engsig-Karup et al. (2008, 2011). The numerical model is implemented in the DTU GPUlab library developed at DTU Compute and uses Nvidia’s CUDA C programming model and is executed on a recent generation of programmable Nvidia GPUs. In performance tests, we demonstrate good and consistent weak scalability of up to 8192 GPUs and on three different compute clusters. We have detailed how the model is capable of solving systems of equations with more than 100 billion degrees of freedom for the first time with the current scalability tests restricted by the amount of hardware available at the time of writing. Also, we highlight that multi-GPU implementations can be a means for further acceleration of run-times in comparison with single GPU computations.

As a first-step, we have demonstrated 2-D curvilinear transformation in the horizontal plane to allow treatment of fixed bottom-mounted structures extending vertically throughout the depth of the fluid. This opens a large class of coastal geometries to a fully nonlinear analysis. An example of wave run-up on a mono-pile structure in open water has been demonstrated with good agreement to experimental results. This approach provides the basis for addressing challenges related to problems of engineering interests: (i) real-time computations, (ii) tsunami propagation over large scales, (iii) marine offshore hydrodynamics, (iv) renewable energy applications, and so on. We restricted our focus to deal with multi-GPU scalability and leave more realistic applications as a part of future work.

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